

```

1 let  $C[0..n-1]$  be a new array;
2 for  $i = 0$  to  $n-1$  do
3   |  $C[i] \leftarrow 0$ ;
4 end
5 for  $i = 0$  to  $n-2$  do
6   for  $j = i+1$  to  $n-1$  do
7     if  $A[i] < A[j]$  then
8       |  $C[j] \leftarrow C[j] + 1$ ;
9     end
10    else
11      |  $C[i] \leftarrow C[i] + 1$ ;
12    end
13  end
14 end
15 for  $i = 0$  to  $n-1$  do
16   |  $B[C[i]] \leftarrow A[i]$ ;
17 end
18 return  $B$ ;

```

Algorithm 14: comparison-counting-sort($A[0..n-1]$)

$$K = h - l$$

```

1 let  $C[0..h-l]$  be a new array;
2 for  $i = 0$  to  $h-l$  do }  $\Theta(K)$ 
3   |  $C[i] \leftarrow 0$ ;
4 end
5 for  $i = 0$  to  $n-1$  do }  $\Theta(n)$ 
6   |  $C[A[i]-l] \leftarrow C[A[i]-l] + 1$ ;
7 end
8 for  $j = 1$  to  $h-l$  do }  $\Theta(K)$ 
9   |  $C[j] \leftarrow C[j] + C[j-1]$ ;
10 end
11 for  $i = n-1$  downto 0 do }  $\Theta(n)$ 
12   |  $j \leftarrow A[i]-l$ ;
13   |  $B[C[j]-1] \leftarrow A[i]$ ;
14   |  $C[j] \leftarrow C[j] - 1$ ;
15 end
16 return  $B$ ;

```

Algorithm 15: counting-sort($A[0..n-1], l, h$)

$$\begin{aligned}
T(n) &= \sum_{i=0}^{n-2} \sum_{j=i+1}^{h-1} 1 = \\
&\sum_{i=0}^{n-2} (n-1-(i+1)+1) = \\
&\sum_{i=0}^{n-2} (n-i-1) = \sum_{i=1}^{n-1} i = \\
&\frac{(n-1) \cdot n}{2} = \Theta(n^2)
\end{aligned}$$

$$\Rightarrow T(n) = \Theta(n+k)$$

Se $K = O(n)$ então

$$\underline{T(n)} = \underline{\Theta(n)}$$